

# On B and Event-B: Principles, Success and Challenges

Jean-Raymond Abrial

[jrabrial@neuf.fr](mailto:jrabrial@neuf.fr)

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- Let us have a **little example**.
- I used this example in an **undergraduate course** recently.
- I know you are **not undergraduate students any more!**
- But it's good sometimes to have some **refreshments!**
- **It's not an example in:** ASM, ALLOY, B, Event-B, TLA, VDM, Z, ...
- **It's an example in:** C. I suppose everyone knows C

```
#define black 0
#define white 1
#define m 5

int color[m] = {black, black, white, black, white}

int print_black_color_index() {
    int k;
    printf("black = ");
    for(k=0; k<=m; k++)
        if (color[k]==white) printf("%d ",k);
    return(0);
}
```

**Is this C program correct?** Let us execute it

```
#define black 0
#define white 1
#define m 5

int color[m] = {black, black, white, black, white}
int print_black_color_index() {
    int k;
    printf("black = ");
    for(k=0; k<=m; k++)
        if (color[k]==white) printf("%d ",k);
    return(0);
}
```

- We obtain the **wrong result**: black = 2 4 (although correct for white)

- The **name** of the program is: "print\_black\_color\_index".
- However, the program prints **white color index instead of black**.
- Does it means that the program is **not correct**?
- Does the **name** of the program say what the **program should do**?

- How to **ensure** that the program is correct?
- We need a statement **outside the program** saying what it should do.
- Now, let us **modify this program** and execute it again

```
#define black 0
#define white 1
#define m 5

int color[m] = {black, black, white, black, white}
int print_white_color_index() {
    int k;
    printf("black = ");
    for(k=0; k<=m; k++)
        if (color[k]==black) printf("%d ",k);
    return(0);
}
```

- We obtain the **strange result**: black = 0 1 3 5

- We just changed **white to black**
- How come that **we get 5**?
- We did **not get any issue with white**
- But now the program acts in a **strange manner with black**
- Let us **look again at the program**
- **Can you guess** what has happened?

```
#define black 0
#define white 1
#define m 5

int color[m] = {black, black, white, black, white}
int print_black_color_index() {
    int k;
    printf("black = ");
    for(k=0; k<=m; k++)
        if (color[k]==black) printf("%d ",k);
    return(0);
}
```

- We obtain the **strange result**: black = 0 1 3 5

- The test in the loop is " $k \leq m$ " although it should be " $k < m$ "
- This had the effect to have the loop going **off the array**
- In fact, the program continues to read the memory **after the array**
- Notice that **C** does **not say anything** in this case

- It happens that **black is 0**
- It happens that the memory after the end of the array **contains a 0**
- As **white is 1**, this error was **not discovered with white**
- If the **memory** would have been **1**, **white** would have been **wrong**
- If the **memory** would have **not been 0 or 1**, the program would have been **"correct" in both cases** (although wrong in fact)
- So, this error discovery is **very arbitrary**

- The C language is **misleading**
- Array indices in C **start at 0**
- So, when we write "**int color[m]**", the **last index value is m-1**
- This is the reason why having "**k<=m**" instead of "**k<m**" led to an error

- But, again, this error was **not discovered** with white, only with black
- The same program was **"correct"** with white and **wrong** with black
- So, **testing is not a reliable approach** to programming
- Does this happen with other programming languages (i.e. **Java**)?
- Yes and no: Java raises an **exception** in this case (and many more)
- Such an exception can be treated by an **exception handler**

- Exception handlers are **terrible features**
- What to do if an **exception** occurs while the **aircraft is about to land**
- But people like them because they are **useful** in the **testing phase**
- However, still dangerous (e.g. **Ariane 5 crash**)
- Still a treatment performed on the **final program** (like **testing**)
- How about **abstract interpretation**?

- The **exception** can be raised (systematically?) **in the laboratory**
- But still a treatment performed on the **final program** (like **testing**)
- **It does not say that the program is correct**
- It just detects that the program does **not contain bad things**:
  - array bound overflow
  - null pointers
  - division by zero
  - ...
- Even **no loop termination checking** (I think)

- What we need is to prove that our program **does not contain bad things** like this
- And many more: **the program meets its specification**
- And many more: the program is **correct within an entire system**
- The answer is in **B** and **Event-B**
- And, clearly, **in other formalisms presented here at ABZ**
- **End of preamble**

- **Basic principles** of B and Event-B
- **Differences** and **similarities** between B and Event-B
- B and Event-B **spreading**
- **Issues** and **challenges**

- "CxC" is the main purpose of B and Event-B
- So, B and Event-B are not programming languages
- They are rather intellectual modelling tool
- They correspond to a practice used in other engineering disciplines

- Refinement is fundamental in engineering practice
- Because a model cannot be defined in a single step
- Models are thus constructed gradually
- From an initial very abstract view to a final concrete one
- Abstraction is usually very difficult to master by informaticians
- So, practitioners have to be seriously educated on this matter

- Refinement is **not sufficient**
- At each step of the development some **statements have to be proved**
- Such statements intend **to ensure that each step is valid**
- For such statements, **no "new language" is developed**
- The **most classical mathematical notation** is used
- Such a notation is that of **predicate logic** and **typed set theory**

- Model cannot be developed **with pen and papers**
- Because they are **too big and too complex**
- Tools are thus **absolutely necessary** to help users writing models
- The tool **for B** is **Atelier B**
- The tool **for Event-B** is **Rodin**
- **Both are free**

- Among the tools, **one is very important**
- **It analyses models** provided by users
- And generate "**proof obligations**" necessary **to validate models**
- This **cannot be done manually** by users as far too much error prone
- Such a "proof obligation generator" was **inspired by that for VDM**

- Once **proof obligations** are generated, they **must be proved**
- For this, **some proving tools** have been **developed** or **imported**
- Such tools work either **automatically** or **interactively**

- **Other tools** are developed
- **By Universities** (Southampton, Düsseldorf, Turku, ...)
- **By industries** (Siemens Transport, Clearsy, Systemrel, ...)
- These tools are **the following** (among others):
  - Model checking
  - Animation
  - Automatic refinement
  - Model decomposition and structuring
  - Link with UML
  - Data validation
  - ...

- Book on B published in 1996
- Book on Event-B published in 2010
- So, Event-B took advantage of this time difference
- Event-B strongly influenced by Action System  
(R-J. Back and R. Kurki-Suonio)
- Event-B is used for the modelling of entire systems

- Main difference: **operations** (in B) and **events** (in Event-B)
- Operations are **pre-conditioned**, whereas events are **guarded**
- Pre-conditions determine **when an operation can be called**
- Guards determine **when an event can occur**
- Both are **assumed in proofs**
- But **pre-conditions are weakened** in refinements
- Whereas **guards are strengthened** in refinements

- Differences in refinements allows **events to be developed gradually**
- **Event parameters** can be **modified, added or instantiated** in refinements
- This is **not possible for operations** (fixed structure from abstraction)
- **New events** can be added in refinements
- **Constants** defined in **separate components** in Event-B
- This allows for **more flexibility** in Event-B than in B

- Event-B has **no programming constructs** as B does (conditionals, loops, sequencing, ...)
- Proof obligations are thus **simpler in Event-B** than in B
- This simplification is important because of the **absence of sequencing in Event-B**
- However, **code generation is simpler for B** than for Event-B

- In Event-B, we worked a lot on **welldefinedness** proof obligations
- Examples:  $f(x)$ ,  $\text{card}(S)$ ,  $\text{min}$ ,  $\text{max}$ ,  $a/b$

- **Usage of the math notation**: predicate logic and typed set theory
- So, proof obligations can be handled by **similar provers in both**
- Consequence: some **provers of Atelier B are used in Rodin**.
- Also both use **similar external provers (SMT provers)**
- **Event-B can be simulated in B** (adding specific proof obligations)

- B is extensively used in Industries (by Cleary and others)
- Cleary claims to make 30% of its business with B
- The main industrial activity is with train systems
- Alstom and Siemens Transport actively participate in these activities
- Train systems with B in Europe, North and South America, Asia

- Event-B widely spread in universities around the world
- In France, United Kingdom, Germany, Spain, Finland
- Also in North and South America (Canada, Brazil, Columbia) and Asia (China, Japan)

- Poor spreading in industries except in train industries
- Other industries could use B (or other formal methods) but do not
- Examples: energy, automotive, aeronautics, space, ...
- People there claim it is too expensive
- They also claim that introducing this technology is difficult
- Event-B and B not used yet in the same project  
(system with Event-B and then software with B)

- For B: [www. clearsy.com/en](http://www.clearsy.com/en)
- For Event-B: [www.Event-B.org](http://www.Event-B.org)
- R.-J. Back and R. Kurki-Suonio *Decentralisation of Process Net with Centralized Control*. Distributed Computing (1989)
- T. Lecomte and al. *Applying a Formal Method in Industry: a 25 years Trajectory*.  
Formal Methods: Foundations and Applications. Springer (2017)